

**Last time:** For a graph  $G = (V, E)$ , we defined:

- (1) The **independence number**  $\alpha(G)$  is the size of a largest independent set in  $G$ .
- (2) The **edge-independence number**  $\alpha'(G)$  is the size of a largest matching in  $G$ .
- (3) A **vertex-cover** of  $G$  is a subset  $U \subset V$  such that every edge of  $G$  has at least one end-point in  $U$ .  
The **vertex-covering number**  $\beta(G)$  is the size of a smallest vertex-cover of  $G$ .
- (4) An **edge-cover** of  $G$  is a subset  $S \subset E$  so that every vertex of  $G$  is incident to at least one edge in  $S$ .  
The **edge-covering number**  $\beta'(G)$  is the size of a smallest edge-cover of  $G$ .

**Example.**

**Today:**

- $\alpha'(G) \leq \beta(G) \leq 2\alpha'(G)$  for any graph.
- $\alpha'(G) = \beta(G)$  if  $G$  is bipartite.
- $\alpha'(G) + \beta'(G) = |V(G)|$  if  $G$  has no isolated vertices.
- $\alpha(G) + \beta(G) = |V(G)|$  for any graph (homework).

**Theorem.** For any graph  $G$  we have  $\alpha'(G) \leq \beta(G) \leq 2\alpha'(G)$ .

**Note:** It is possible that  $\beta(G) = 2\alpha'(G)$ , just take  $G = K_n$  with  $n$  odd.

**Next:**  $\alpha'(G) = \beta(G)$  is also possible.

**Theorem** (König's Theorem). *If  $G$  is a bipartite graph, then  $\alpha'(G) = \beta(G)$ . In other words, for bipartite graphs, the size of a maximum matching is equal to the size of a smallest vertex-cover.*

**Lemma.** *Let  $G = (V, E)$  be a graph with no isolated vertices. If  $S \subset E$  is a minimum edge-cover of  $G$ , then  $(V, S)$  is a forest with no isolated vertices.*

**Theorem (8.7).** *If  $G$  is a graph on  $n$  vertices none of which is isolated, then  $\alpha'(G) + \beta'(G) = n$ .*

**Theorem (8.8).** *If  $G$  is a graph on  $n$  vertices, then  $\alpha(G) + \beta(G) = n$ .*

**Example.** *Compute  $\alpha(P)$  and  $\beta(P)$  where  $P$  is the Petersen graph.*

**Next:** a nice consequence of Hall's theorem.

**Theorem.** *Let  $G$  be a bipartite graph with parts  $U, W$  such that  $|U| \leq |W|$ . Assume that  $G$  has no isolated vertices and that  $\deg u \geq \deg w$  whenever  $u \in U$  and  $w \in W$  are such that  $uw \in E(G)$ . Then  $G$  has a matching which saturates  $U$ .*